

Augmented Marked Graphs

King-Sing Cheung
 University of Hong Kong
 Pokfulam, Hong Kong
 E-mail: ks.cheung@hku.hk

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Augmented marked graphs possess some structural characteristics desirable for modelling shared resource systems such as manufacturing systems. However, there are only a few known properties on augmented marked graphs, and these known properties are mainly on liveness and reversibility. In this paper, the properties of augmented marked graphs are reviewed extensively. Siphon-based and cycle-based characterisations for liveness and reversibility as well as transformation-based characterisations for boundedness and conservativeness are proposed. Pretty simple conditions and procedures are then derived for checking the liveness, reversibility, boundedness and conservativeness of augmented marked graphs. The dining philosopher problem is used for illustration.

Povzetek: Opisane so lastnosti grafov za predstavitev sistemov z deljenimi viri.

1 Introduction

Augmented marked graphs were first introduced by Chu and Xie [1]. They are not well known as compared to other sub-classes of Petri nets such as free-choice nets [2], and the properties of augmented marked graphs are not studied extensively. However, augmented marked graphs possess a structure which is desirable for modelling shared resources, and for this reason, they are often used in modelling shared resource systems, such as manufacturing systems [1, 3, 4, 5, 6, 7].

In the literature, the studies of augmented marked graphs mainly focus on deadlock-freeness, liveness and reversibility. Based on mathematical programming, Chu and Xie proposed a necessary and sufficient condition of live and reversible augmented marked graphs, which checks the existence of potential deadlocks [1]. However, this involves analysis on the flow of tokens during execution and the checking cannot be simply made by looking into the structure. Chu and Xie also proposed a siphon-based characterisation for live and reversible augmented marked graphs but it provides a sufficient condition only. The boundedness and conservativeness of augmented marked graphs were not investigated.

There are other studies of augmented marked graphs, which are mainly on the property-preserving synthesis or composition of augmented marked graphs. Jeng proposed a synthesis method of process nets for manufacturing system design [4, 5]. (Note : Process nets broadly cover augmented marked graphs.) Based on siphons and the firability of transitions, sufficient conditions for liveness and reversibility are derived. Huang also investigated the composition of augmented marked graphs via common resource places, so that some essential properties such as liveness, boundedness and reversibility can be preserved under certain conditions [6].

In our previous works on augmented marked graphs, we proposed new characterisations for live and reversible augmented marked graphs as well as the synthesis of augmented marked graphs for system design [7, 8, 9, 10, 11]. This paper extends our previous works with a focus on the properties of augmented marked graphs. It reports the following two contributions.

First, we propose a number of characterisations for live and reversible augmented marked graphs, based on siphons and cycles. In particular, a new property called R-inclusion property is introduced to characterise the siphon-trap property of augmented marked graphs. With this property, a pretty simple necessary and sufficient condition for live and reversible augmented marked graphs is then proposed. Second, for analysis of the boundedness and conservativeness of augmented marked graphs, a R-transform is introduced to transform an augmented marked graph into marked graphs. With the R-transform, a pretty simple necessary and sufficient condition for bounded and conservative augmented marked graphs is proposed. These characterisations will be illustrated using the dining philosopher problem.

The rest of this paper is organised as follows. Following this introduction, Section 2 provides the preliminaries to be used in this paper. Section 3 briefly introduces augmented marked graphs. Section 4 focus on liveness and reversibility of augmented marked graphs, where siphon-based and cycle-based characterisations are proposed. Section 5 then focus on boundedness and conservativeness of augmented marked graphs, where transformation-based characterisations are proposed. Section 6 illustrates these characterisations using the dining philosopher example. Finally, Section 7 concludes our results.

It should be noted that, in this paper, proofs of the proposed properties are shown in the appendix.

2 Preliminary

This section provides the preliminaries to be used in this paper for those readers who are not familiar with Petri nets [12, 13, 14, 15].

A place-transition net (PT-net) is a directed graph consisting of two sorts of nodes called places and transitions, such that no arcs connect two nodes of the same sort. Graphically, a place is denoted by a circle, a transition by a box, and an arc by a directed line. A Petri net is a PT-net with tokens assigned to its places, and the token distribution is denoted by a marking.

A Petri net is usually used to represent a discrete system, where the places denote conditions, the transitions denote events and the arcs between places and transitions denote the relationship between conditions and events.

Definition 1. A place-transition net (PT-net) is a 4-tuple $N = \langle P, T, F, W \rangle$, where P is a set of places, T is a set of transitions, $F \subseteq (P \times T) \cup (T \times P)$ is a flow relation and $W : F \rightarrow \{ 1, 2, \dots \}$ is a weight function. N is said to be ordinary if and only if the range of W is $\{ 1 \}$.

An ordinary PT-net is usually written as $\langle P, T, F \rangle$. In the rest of this paper, unless specified otherwise, all PT-nets refer to ordinary PT-nets.

Definition 2. Let $N = \langle P, T, F, W \rangle$ be a PT-net. For any $x \in (P \cup T)$, $\bullet x = \{ y \mid (y, x) \in F \}$ and $x^\bullet = \{ y \mid (x, y) \in F \}$ are called the pre-set and post-set of x , respectively.

For clarity in presentation, the pre-set and post-set of a set of places or transitions $X = \{ x_1, x_2, \dots, x_n \}$ can be written as $\bullet X$ and X^\bullet respectively, where $\bullet X = \bullet x_1 \cup \bullet x_2 \cup \dots \cup \bullet x_n$ and $X^\bullet = x_1^\bullet \cup x_2^\bullet \cup \dots \cup x_n^\bullet$.

Definition 3. For a PT-net $N = \langle P, T, F, W \rangle$, a path is a sequence of places and transitions $\rho = \langle x_1, x_2, \dots, x_n \rangle$, such that $(x_i, x_{i+1}) \in F$ for $i = 1, 2, \dots, n-1$. ρ is said to be elementary if and only if it contains no duplicate places or transitions.

Definition 4. For a PT-net $N = \langle P, T, F, W \rangle$, a sequence of places $\langle p_1, p_2, \dots, p_n \rangle$ is called a cycle if and only if there exists a set of transitions $\{ t_1, t_2, \dots, t_n \}$, such that $\langle p_1, t_1, p_2, t_2, \dots, p_n, t_n \rangle$ forms an elementary path and $(t_n, p_1) \in F$.

Definition 5. For a PT-net $N = \langle P, T, F, W \rangle$, a marking is a function $M : P \rightarrow \{ 0, 1, 2, \dots \}$, where $M(p)$ is the number of tokens in p . (N, M_0) represents N with an initial marking M_0 .

Semantically, a marking represents the state of a Petri net. The initial marking specifically represents the initial state of a Petri net. A transition is enabled and can be fired at a state (marking) where all the places in its pre-set hold tokens. On firing the transition, tokens will be moved from the places in its pre-set to the places in its post-set. The firing of a transition is formally defined as follows.

Definition 6. For a PT-net (N, M_0) , a transition t is said to be enabled at a marking M if and only if $\forall p \in \bullet t : M(p) \geq W(p, t)$. On firing t , M is changed to M' such that $\forall p \in P : M'(p) = M(p) - W(p, t) + W(t, p)$. In notation, $M [N, t] M'$ or $M [t] M'$.

Definition 7. For a PT-net (N, M_0) , a sequence of transitions $\sigma = \langle t_1, t_2, \dots, t_n \rangle$ is called a firing sequence if and only if $M_0 [t_1] \dots [t_n] M_n$. In notation, $M_0 [N, \sigma] M_n$ or $M_0 [\sigma] M_n$.

Definition 8. For a PT-net (N, M_0) , a marking M is said to be reachable if and only if there exists a firing sequence σ such that $M_0 [\sigma] M$. In notation, $M_0 [N, *] M$ or $M_0 [*] M$. $[N, M_0]$ or $[M_0]$ represents the set of all reachable markings of (N, M_0) .

The structure of a PT-net can be represented by a matrix called incidence matrix.

Definition 9. Let $N = \langle P, T, F, W \rangle$ be a PT-net, where $P = \{ p_1, p_2, \dots, p_m \}$ and $T = \{ t_1, t_2, \dots, t_n \}$. The incidence matrix of N is an $m \times n$ matrix V whose typical entry $v_{ij} = W(p_i, t_j) - W(t_j, p_i)$ represents the change in number of tokens in p_i after firing t_j once, for $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$.

Liveness, boundedness, safeness, reversibility and conservativeness are best known properties of Petri nets. Liveness implies freeness of deadlocks. Boundedness and safeness imply freeness of capacity overflow. Reversibility refers to the capability of being reinitialised from any reachable states. Conservativeness is a special form of boundedness.

Definition 10. For a PT-net (N, M_0) , a transition t is said to be live if and only if $\forall M \in [M_0], \exists M' : M [*] M' [t]$. (N, M_0) is said to be live if and only if every transition is live.

Definition 11. For a PT-net (N, M_0) , a place p is said to be k -bounded if and only if $\forall M \in [M_0] : M(p) \leq k$, where k is a positive integer. (N, M_0) is said to be bounded if and only if every place is k -bounded, and safe if and only if every place is 1-bounded.

Definition 12. A PT-net (N, M_0) is said to be reversible if and only if $\forall M \in [M_0] : M [*] M_0$.

Definition 13. For a PT-net $N = \langle P, T, F, W \rangle$, a place invariant is a $|P|$ -vector $\alpha \geq 0$ such that $\alpha V = 0$, where V is the incidence matrix of N .

Definition 14. A PT-net is said to be conservative if and only if there exists a place invariant $\alpha > 0$.

Figure 1 shows an ordinary PT-net which is live, bounded, safe, reversible and conservative.

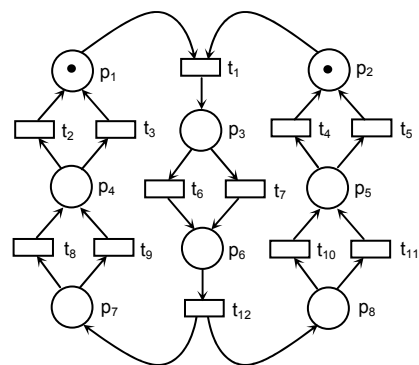


Figure 1: A live, bounded, safe, reversible and conservative PT-net.

Property 1. A PT-net (N, M_0) is bounded if it is conservative [14, 15].

Definition 15. For a PT-net N , a set of places S is called a siphon if and only if $\bullet S \subseteq S^*$. S is said to be minimal if and only if there does not exist another siphon S' in N such that $S' \subset S$.

Definition 16. For a PT-net, a set of places T is called a trap if and only if $T^* \subseteq \bullet T$.

Definition 17. A PT-net (N, M_0) is said to satisfy the siphon-trap property if and only if every siphon contains a marked trap (or every minimal siphon contains a marked trap).

A well known sub-class of Petri nets, marked graphs possess many special properties pertaining to its liveness, boundedness and reversibility.

Definition 18. A marked graph is an ordinary PT-net $N = \langle P, T, F, W \rangle$ such that $\forall p \in P : |\bullet p| = |p^*| = 1$.

Property 2. A marked graph (N, M_0) is live if and only if every cycle is marked by M_0 [13, 14].

Property 3. A live marked graph (N, M_0) is bounded if and only if every place belongs to a cycle marked by M_0 [13, 14].

Property 4. A live and bounded marked graph is reversible [13, 14].

Property 5. For a marked graph, the corresponding place vector of a cycle is a place invariant [13, 14].

Figure 2 shows a marked graph which is live, bounded, safe and reversible. Places $\langle p_1, p_3, p_6, p_7, p_4 \rangle$ form a cycle. The place vector is a place invariant.

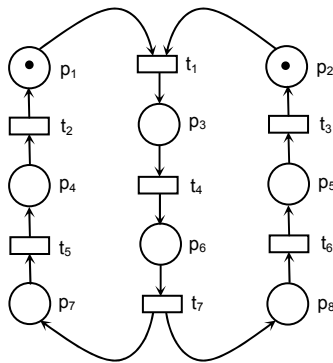


Figure 2: A live, bounded, safe and reversible marked graph.

3 Augmented marked graphs

Augmented marked graphs were first introduced by Chu and Xie [1]. This section briefly describes augmented marked graphs.

Definition 19. An augmented marked graph $(N, M_0; R)$ is a PT-net (N, M_0) with a specific subset of places R , such that : (a) Every place in R is marked by M_0 . (b) The net (N', M_0') obtained from $(N, M_0; R)$ by removing the places in R and their associated arcs is a marked graph. (c) For each $r \in R$, there exist $k_r \geq 1$ pairs of transitions $D_r = \{ \langle t_{s1}, t_{h1} \rangle, \langle t_{s2}, t_{h2} \rangle, \dots, \langle t_{skr}, t_{hkr} \rangle \}$, such that $r^* = \{ t_{s1}, t_{s2}, \dots, t_{skr} \} \subseteq T$ and $\bullet r = \{ t_{h1}, t_{h2}, \dots, t_{hkr} \} \subseteq T$ and that, for each $\langle t_{si}, t_{hi} \rangle \in D_r$, there exists in (N', M_0') an elementary path ρ_{ri} connecting t_{si} to t_{hi} . (d) In (N', M_0') , every cycle is marked and no ρ_{ri} is marked.

Figure 3 shows a typical augmented marked graph $(N, M_0; R)$, where $R = \{ r_1, r_2 \}$. For r_1 , $D_{r1} = \{ \langle t_1, t_{11} \rangle, \langle t_3, t_9 \rangle \}$. For r_2 , $D_{r2} = \{ \langle t_2, t_{11} \rangle, \langle t_4, t_{10} \rangle \}$.

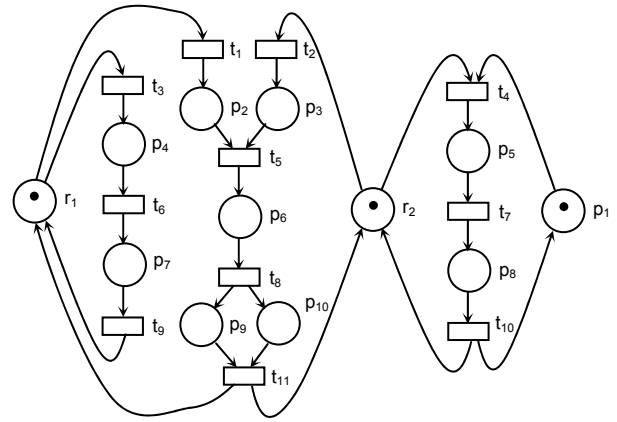


Figure 3: An augmented marked graph.

Augmented marked graphs possess a number of special properties pertaining to liveness, boundedness, reversibility and conservativeness. In the following sections, these properties are thoroughly investigated.

4 Liveness and reversibility

This section focuses on the liveness and reversibility of augmented marked graphs. After reporting several known properties, some siphon-based and cycle-based characterisations for live and reversible augmented marked graphs are proposed.

Property 6. An augmented marked graph is live if and only if it does not contain any potential deadlock [1]. (Note : A potential deadlock is a siphon which would eventually become empty.)

Property 7. An augmented marked graph is reversible if it is live [1].

Property 8. An augmented marked graph is live and reversible if and only if every minimal siphon would never become empty.

Property 9. An augmented marked graph $(N, M_0; R)$ is live and reversible if every minimal siphon, which contains at least one place of R , contains a marked trap [1].

For the augmented marked graph $(N, M_0; R)$ shown in Figure 3, the minimal siphons are : $\{ p_1, p_5, p_8 \}$, $\{ r_1, p_2, p_4, p_6, p_7, p_9 \}$, $\{ r_1, p_2, p_4, p_6, p_7, p_{10} \}$, $\{ r_2, p_3, p_5, p_6, p_8, p_9 \}$ and $\{ r_2, p_3, p_5, p_6, p_8, p_{10} \}$. Each of these minimal siphons contains a marked trap, and would never become empty. $(N, M_0; R)$ is live and reversible.

The places and transitions generated by cycles are defined as follows.

Definition 20. For a PT-net N , Ω_N is defined as the set of all cycles in N .

Definition 21. Let N be a PT-net. For a subset of cycles $Y \subseteq \Omega_N$, $P[Y]$ is defined as the set of places in Y , and $T[Y] = \bullet P[Y] \cap P[Y]^*$ is defined as the set of transitions generated by Y .

For clarity in presentation, $P[\{\gamma\}]$ and $T[\{\gamma\}]$ can be written as $P[\gamma]$ and $T[\gamma]$, to denote the set of places in a cycle γ and the set of transitions generated by γ , respectively.

Definition 22. For a PT-net N , an elementary path $\rho = \langle x_1, x_2, \dots, x_n \rangle$ is said to be conflict-free if and only if, for any transition x_i in ρ , $j \neq (i-1) \Rightarrow x_j \notin \bullet x_i$.

Property 10. Let S be a minimal siphon of a PT-net. For any $p, p' \in S$, there exists in S a conflict-free path from p to p' [16].

Property 11. For a minimal siphon S of an augmented marked graph $(N, M_0; R)$, there exists a set of cycles $Y \subseteq \Omega_N$ such that $P[Y] = S$.

Property 12. Every cycle in an augmented marked graph is marked.

Property 13. Every siphon in an augmented marked graph is marked.

Property 14. Let $(N, M_0; R)$ be an augmented marked graph. For every $r \in R$, there exists a minimal siphon which contains only one marked place r .

Consider the augmented marked graph $(N, M_0; R)$ shown in Figure 3. Every minimal siphon is covered by cycles. Consider a minimal siphon $S_1 = \{ r_1, p_2, p_4, p_6, p_7, p_9 \}$. There exists a set of cycles $Y_1 = \{ \gamma_{11}, \gamma_{12} \}$, where $\gamma_{11} = \langle r_1, p_4, p_7 \rangle$ and $\gamma_{12} = \langle r_1, p_2, p_6, p_9 \rangle$, such that $S_1 = P[Y_1]$. Consider another minimal siphon $S_2 = \{ r_2, p_3, p_5, p_6, p_8, p_{10} \}$. There exists a set of cycles $Y_2 = \{ \gamma_{21}, \gamma_{22} \}$, where $\gamma_{21} = \langle r_2, p_5, p_8 \rangle$ and $\gamma_{22} = \langle r_2, p_3, p_6, p_{10} \rangle$, such that $S_2 = P[Y_2]$. For S_1 , $r_1 \in R$ is the only one marked place. Also, for S_2 , $r_2 \in R$ is the only one marked place.

For an augmented marked graph, minimal siphons can be classified into R-siphons and NR-siphons. Based on R-siphons and NR-siphons, some characterisations for augmented marked graphs are proposed.

Definition 23. For an augmented marked graph $(N, M_0; R)$, a minimal siphon is called a R-siphon if and only if it contains at least one place in R .

Definition 24. For an augmented marked graph $(N, M_0; R)$, a minimal siphon is called a NR-siphon if and only if it does not contain any place in R .

Definition 25. Let N be a PT-net. For a set of places Q in N , $\Omega_N[Q]$ is defined as the set of cycles that contains at least one place in Q .

For clarity in presentation, $\Omega_N[\{p\}]$ can be written as $\Omega_N[p]$ to denote the set of cycles that contains a place p .

Property 15. For an augmented marked graph $(N, M_0; R)$, a R-siphon is covered by a set of cycles $Y \subseteq \Omega_N[R]$.

Figure 4 shows another augmented marked graph $(N, M_0; R)$, where $R = \{ r_1, r_2 \}$. There are five minimal siphons, namely, $S_1 = \{ r_1, p_3, p_4, p_7, p_8 \}$, $S_2 = \{ r_1, p_3, p_5, p_7, p_8 \}$, $S_3 = \{ r_2, p_2, p_4, p_6, p_8, p_9, p_{10} \}$, $S_4 = \{ r_2, p_2, p_5, p_6, p_8, p_9, p_{10} \}$ and $S_5 = \{ p_1, p_3, p_7 \}$. S_1, S_2, S_3 and S_4 are R-siphon as they contain at least one place in R . S_5 is a NR-siphon which does not contain any place in R . For $(N, M_0; R)$, every R-siphon is covered by a set of cycles in $\Omega_N[R]$. For example, $S_1 = \{ r_1, p_3, p_4, p_7, p_8 \}$ is covered by a set of cycles $Y_1 = \{ \gamma_{11}, \gamma_{12} \} \subseteq \Omega_N[R]$, where $\gamma_{11} = \langle r_1, p_3, p_7 \rangle$ and $\gamma_{12} = \langle r_1, p_4, p_8 \rangle$.

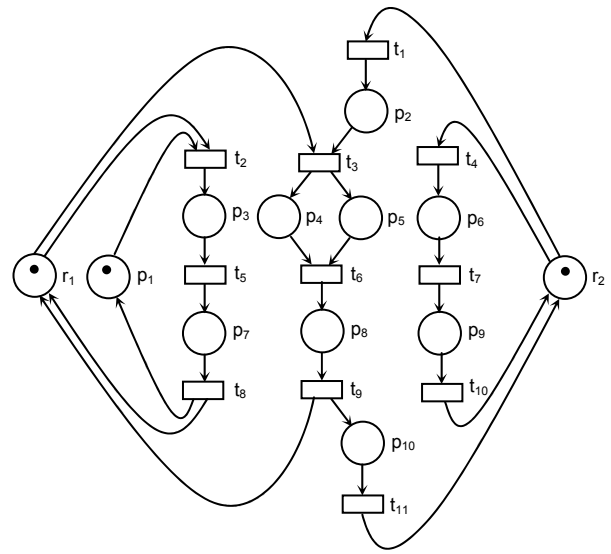


Figure 4: Another augmented marked graph.

Property 16. Let S be a R-siphon of an augmented marked graph $(N, M_0; R)$. For every $t \in (S^* \setminus \bullet S)$, there does not exist any $s \in (S \setminus R)$ such that $t \in s^*$.

Property 17. For an augmented marked graph $(N, M_0; R)$, a NR-siphon contains itself as a marked trap and would never become empty.

Property 18. An augmented marked graph $(N, M_0; R)$ is live and reversible if and only if no R-siphons eventually become empty.

Property 19. An augmented marked graph $(N, M_0; R)$ satisfies the siphon-trap property if and only if every R-siphon contains a marked trap.

Consider the augmented marked graph $(N, M_0; R)$ shown in Figure 4. Every R-siphon contains a marked trap. Each of the R-siphons $S_1 = \{ r_1, p_3, p_4, p_7, p_8 \}$, $S_2 = \{ r_1, p_3, p_5, p_7, p_8 \}$, $S_3 = \{ r_2, p_2, p_4, p_6, p_8, p_9, p_{10} \}$ and $S_4 = \{ r_2, p_2, p_5, p_6, p_8, p_9, p_{10} \}$ contains a marked trap and would never become empty. $(N, M_0; R)$ is live and reversible.

Property 20. (characterisation of Property 9) An augmented marked graph $(N, M_0; R)$ is live and reversible if every R-siphon contains a marked trap.

Property 18 provides a simple necessary and sufficient condition for live and reversible augmented marked graphs. With Properties 18 and 20, we can determine if an augmented marked graph is live and reversible based on R-siphons. Besides, Property 15 provides a characterisation for R-siphons so that R-siphons can be identified by finding cycles in $\Omega_N[R]$. We may now derive a strategy for checking the liveness and reversibility of an augmented marked graph $(N, M_0; R)$:

- (a) Find all R-siphons based on $\Omega_N[R]$.
- (b) Check if every R-siphon contains a marked trap. If yes, report $(N, M_0; R)$ is live and reversible. Otherwise, go to (c).
- (c) For each R-siphon which does not contain any marked trap, check if it would never become empty. If yes, report $(N, M_0; R)$ is live and reversible. Otherwise, report $(N, M_0; R)$ is neither live nor reversible.

In the following, conflict-free cycles are introduced. Based on conflict-free cycles, a new property called R-inclusion is proposed. It is then used for characterising liveness and reversibility of augmented marked graphs.

Definition 26. For a PT-net N , a set of cycles $Y \subseteq \Omega_N$ is said to be conflict-free if and only if, for any $q, q' \in P[Y]$, there exists in $P[Y]$ a conflict-free path from q to q' .

Figure 5 shows a PT-net N . Consider three cycles $\gamma_1, \gamma_2, \gamma_3 \in \Omega_N[p_3]$, where $\gamma_1 = \langle p_3, p_2, p_7 \rangle$, $\gamma_2 = \langle p_3, p_4 \rangle$ and $\gamma_3 = \langle p_3, p_1, p_6, p_{10}, p_8 \rangle$. The set of cycles $Y_1 = \{ \gamma_1, \gamma_2 \}$ is conflict-free because for any $q, q' \in P[Y_1]$, there exists in $P[Y_1]$ a conflict-free path from q to q' . The set of cycles $Y_2 = \{ \gamma_2, \gamma_3 \}$ is not conflict-free. We have $p_4, p_8 \in P[Y_2]$. p_4 is connected to p_8 via only one path $\rho = \langle p_4, t_5, p_3, t_1, p_1, t_3, p_6, t_6, p_{10}, t_9, p_8 \rangle$ in Y_2 , and ρ is not conflict-free because $p_4, p_8 \in \bullet t_5$.

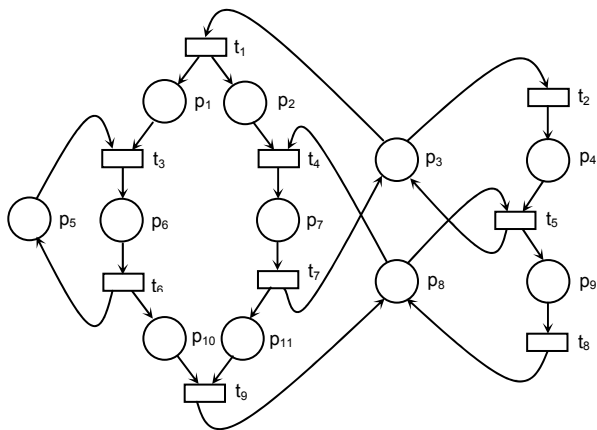


Figure 5: A PT-net for illustration of conflict-free cycles.

Property 21. Let S be a minimal siphon of an augmented marked graph $(N, M_0; R)$, and $Y \subseteq \Omega_N$ be a set of cycles such that $S = P[Y]$. Then, Y is conflict-free.

For the augmented marked graph shown in Figure 3, $\{ r_1, p_2, p_4, p_6, p_7, p_9 \}$ is a minimal siphon covered by a set of cycles $\{ \langle r_1, p_4, p_7 \rangle, \langle r_1, p_2, p_6, p_9 \rangle \}$ which is conflict free. $\{ r_2, p_3, p_5, p_6, p_8, p_{10} \}$ is another minimal siphon covered by a set of cycles $\{ \langle r_2, p_5, p_8 \rangle, \langle r_2, p_3, p_6, p_{10} \rangle \}$ which is conflict-free. For the augmented marked graph shown in Figure 4, $\{ r_1, p_3, p_4, p_7, p_8 \}$ is a minimal siphon covered by a set of cycles $\{ \langle r_1, p_3, p_7 \rangle, \langle r_1, p_4, p_8 \rangle \}$ which is conflict free. $\{ r_1, p_3, p_5, p_7, p_8 \}$ is another minimal siphon covered by a set of cycles $\{ \langle r_1, p_3, p_7 \rangle, \langle r_1, p_5, p_8 \rangle \}$ which is conflict free.

Definition 27. Let $(N, M_0; R)$ be an augmented marked graph. A place $r \in R$ is said to satisfy the R-inclusion if and only if, for any set of cycles $Y \subseteq \Omega_N[R]$ such that Y is conflict-free, $\bullet r \subseteq T[Y] \Rightarrow r^\bullet \subseteq T[Y]$.

Figure 6 shows an augmented marked graph $(N, M_0; R)$, where $R = \{ r_1, r_2 \}$. Consider r_1 . For any set of cycles $Y_1 \subseteq \Omega_N[R]$ such that Y_1 is conflict-free, $\bullet r_1 \subseteq T[Y_1] \Rightarrow r_1^\bullet \subseteq T[Y_1]$. Next, consider r_2 . For any set of cycles $Y_2 \subseteq \Omega_N[R]$ such that Y_2 is conflict-free, $\bullet r_2 \subseteq T[Y_2] \Rightarrow r_2^\bullet \subseteq T[Y_2]$. Both r_1 and r_2 satisfy the R-inclusion.

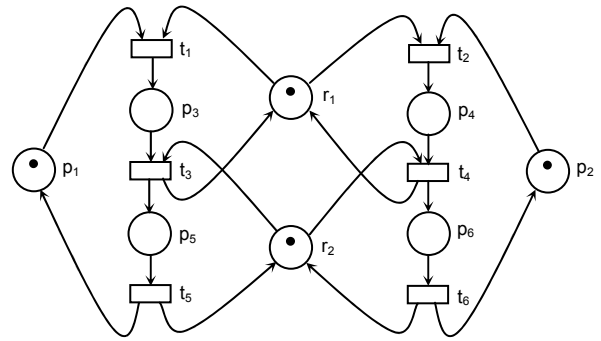


Figure 6: An augmented marked graph for illustration of R-inclusion.

Figure 7 shows another augmented marked graph $(N, M_0; R)$. For $r_1 \in R$, there exists a set of cycles $Y_1 = \{ \gamma_{11}, \gamma_{12} \} \subseteq \Omega_N[R]$, where $\gamma_{11} = \langle r_1, p_5 \rangle$ and $\gamma_{12} = \langle r_1, p_5, r_2, p_6 \rangle$. $\bullet r_1 = \{ t_5, t_6 \} \subseteq T[Y_1] = \{ t_3, t_4, t_5, t_6 \}$ but $r_1^\bullet = \{ t_2, t_3 \} \not\subseteq T[Y_1]$. For $r_2 \in R$, there exists a set of cycles $Y_2 = \{ \gamma_{21}, \gamma_{22} \} \subseteq \Omega_N[R]$, where $\gamma_{21} = \langle r_2, p_6 \rangle$ and $\gamma_{22} = \langle r_2, p_6, r_1, p_5 \rangle$. $\bullet r_2 = \{ t_5, t_6 \} \subseteq T[Y_2] = \{ t_3, t_4, t_5, t_6 \}$ but $r_2^\bullet = \{ t_1, t_4 \} \not\subseteq T[Y_2]$. Both r_1 and r_2 do not satisfy the R-inclusion property.

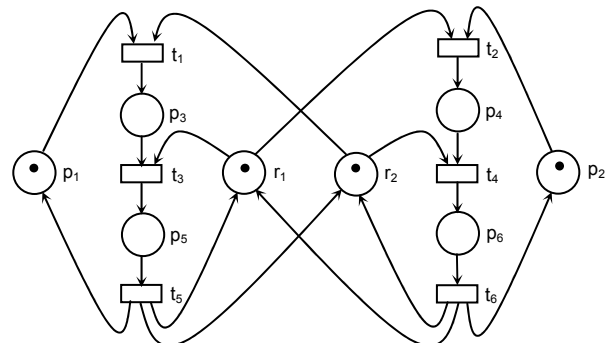


Figure 7: Another augmented marked graph for illustration of R-inclusion.

Property 22. For an augmented marked graph $(N, M_0; R)$, a R-siphon S contains itself as a marked trap if every place $r \in R$ in S satisfies the R-inclusion property.

Property 23. An augmented marked graph $(N, M_0; R)$ satisfies the siphon-trap property if and only if every place $r \in R$ satisfies the R-inclusion property.

Property 24. An augmented marked graph $(N, M_0; R)$ is live and reversible if every place $r \in R$ satisfies the R-inclusion property.

Consider the augmented marked graph $(N, M_0; R)$ shown in Figure 6. We have $R = \{ r_1, r_2 \}$, where both r_1 and r_2 satisfy the R-inclusion property. Any R-siphon, such as $\{ r_1, p_3, p_4 \}$ or $\{ r_2, p_5, p_6 \}$, contains itself as a marked trap. $(N, M_0; R)$ satisfies the siphon-trap property, and is live and reversible.

Property 24 provides a cycle-based condition for live and reversible augmented marked graphs. We need to check the R-inclusion property which involves finding cycles and checking their pre-sets and post-sets.

Based on Properties 15, 18, 20, 22 and 24, we revise the strategy for checking the liveness and reversibility of an augmented marked graph $(N, M_0; R)$ with the use of the R-inclusion property, as follows.

(a) Check if every $r \in R$ satisfies the R-inclusion property. If yes, report $(N, M_0; R)$ is live and reversible. Otherwise go to (b).

(b) Let $R' \subseteq R$ be the set of places which do not satisfy the R-inclusion property. Based on $\Omega_{2N}[R']$, find all R-siphons which contain at least one place in R' .

(c) For each R-siphon identified in (b), check if it contains a marked trap. If yes, report $(N, M_0; R)$ is live and reversible. Otherwise, go to (d).

(d) For each R-siphon identified in (b) that does not contain any marked trap, check if it would never become empty. If yes, report $(N, M_0; R)$ is live and reversible. Otherwise, report $(N, M_0; R)$ is neither live nor reversible.

5 Boundedness and conservativeness

This section focus on the boundedness and conservativeness of augmented marked graphs, which are less studied in the literature. Some transform-based characterisations for bounded and conservative augmented marked graphs are proposed.

In the following, we introduce a new transformation called R-transform for augmented marked graphs. It simply transforms an augmented marked graphs $(N, M_0; R)$ into a number of marked graphs by replacing each place in R by a set of places.

Property 25. Let $(N, M_0; R)$ be an augmented marked graph to be transformed into (N', M_0') as follows. For each place $r \in R$, where $D_r = \{ \langle t_{s1}, t_{h1} \rangle, \langle t_{s2}, t_{h2} \rangle, \dots, \langle t_{skr}, t_{hkr} \rangle \}$, replace r with a set of places $\{ p_1, p_2, \dots, p_{kr} \}$, such that $M_0'[p_i] = M_0[r]$ and $p_i \bullet = \{ t_{si} \}$ and $p_i^\bullet = \{ t_{hi} \}$ for $i = 1, 2, \dots, k_r$. Then, (N', M_0') is a marked graph.

Definition 28. Let $(N, M_0; R)$ be an augmented marked graph. The marked graph (N', M_0') transformed from $(N, M_0; R)$ as stated in Property 25 is called the R-transform of $(N, M_0; R)$.

Property 26. The R-transform of an augmented marked graph is live.

Figure 8 shows an augmented marked graph. Figure 9 shows its R-transform, where r is replaced by $\{ q_1, q_2 \}$.

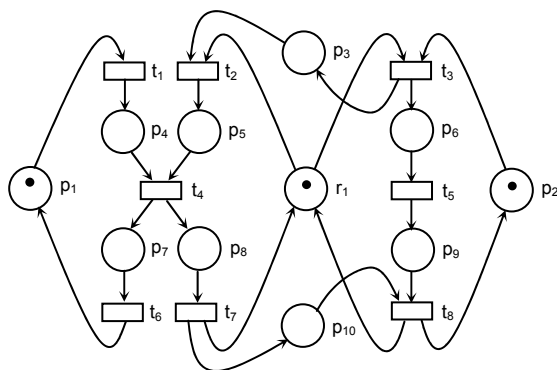


Figure 8: An augmented marked graph for illustration of R-transform.

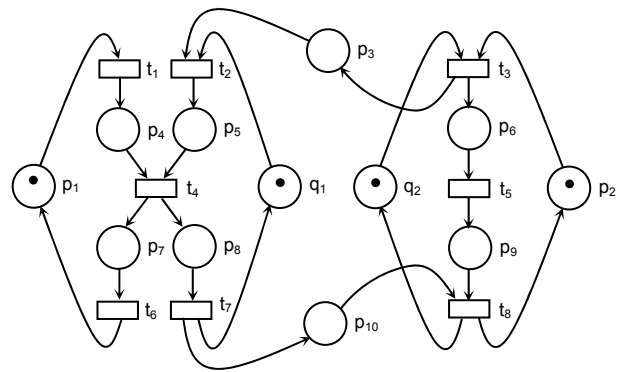


Figure 9: The R-transform of the augmented marked graph shown in figure 8.

Property 27. Let (N', M_0') be the R-transform of an augmented marked graph $(N, M_0; R)$, where $r \in R$ is replaced by a set of places $Q = \{ q_1, q_2, \dots, q_k \}$, and P_0 be the set of marked places in N' . Then, for each q_i in N' , there exists a place invariant α_i of N' such that $\alpha_i[q_i] = 1$ and $\alpha_i[s] = 0$ for any place $s \in P_0 \setminus \{q_i\}$.

Property 28. Let $(N, M_0; R)$ be an augmented marked graph, where $R = \{ r_1, r_2, \dots, r_n \}$. Let (N', M_0') be the R-transform of $(N, M_0; R)$, where each r_i is replaced by a set of places Q_i , for $i = 1, 2, \dots, n$. If every place in (N', M_0') belongs to a cycle, then there exists a place invariant α of N' such that $\alpha > 0$ and $\alpha[q_1] = \alpha[q_2] = \dots = \alpha[q_k]$ for each $Q_i = \{ q_1, q_2, \dots, q_k \}$.

Consider the R-transform (N', M_0') shown in Figure 9. It is a live marked graph. For q_1 , there exists a place invariant α_1 , such that $\alpha_1[q_1] = 1$ and $\alpha_1[q_2] = \alpha_1[p_1] = \alpha_1[p_2] = 0$. For q_2 , there also exists a place invariant α_2 , such that $\alpha_2[q_2] = 1$ and $\alpha_2[q_1] = \alpha_2[p_1] = \alpha_2[p_2] = 0$. In (N', M_0') , every place belongs to a cycle. There also exists a place invariant $\alpha > 0$, where $\alpha[q_1] = \alpha[q_2]$.

Based on R-transform, a necessary and sufficient condition for bounded and conservative augmented marked graphs is proposed.

Property 29. Let (N', M_0') be the R-transform of an augmented marked graph $(N, M_0; R)$. $(N, M_0; R)$ is bounded and conservative if and only if every place in (N', M_0') belongs to a cycle.

With Properties 29, we derive the following strategy for checking the boundedness and conservativeness of an augmented marked graph $(N, M_0; R)$:

- (a) Create the R-transform (N', M_0') of $(N, M_0; R)$.
- (b) For each place p in (N', M_0') , check if there exists a cycle that contains p . If yes, report $(N, M_0; R)$ is bounded and conservative. Otherwise, report $(N, M_0; R)$ is neither bounded nor conservative.

Property 30. Let (N', M_0') be the R-transform of an augmented marked graph $(N, M_0; R)$. $(N, M_0; R)$ is bounded and conservative if and only if (N', M_0') is bounded.

Consider the augmented marked graph $(N, M_0; R)$ shown in Figure 8, and the R-transform (N', M_0') of $(N, M_0; R)$ in Figure 9. Every place in (N', M_0') belongs to a cycle. $(N, M_0; R)$ is bounded and conservative. (N', M_0') is also bounded and conservative.

Figure 10 shows an augmented marked graph $(N, M_0; R)$, and Figure 11 shows the R-transform (N', M'_0) of $(N, M_0; R)$. For (N', M'_0) , p_3 does not belong to any cycle. Also, p_8 does not belong to any cycle. $(N, M_0; R)$ is neither bounded nor conservative.

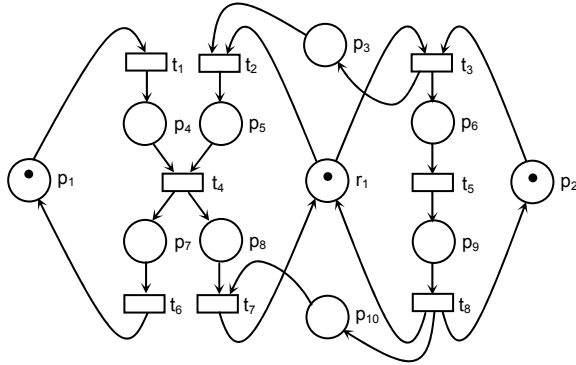


Figure 10: Another augmented marked graph for illustration of R-transform.

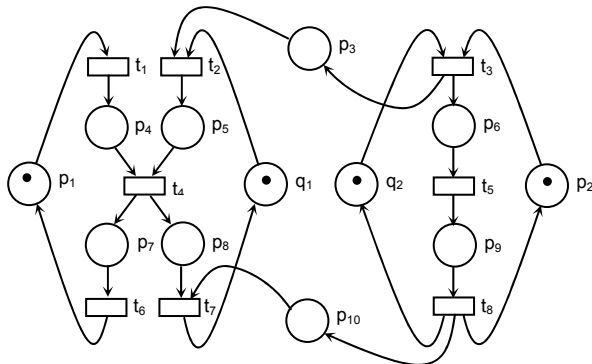


Figure 11: The R-transform of the augmented marked graph shown in figure 10.

6 The dining philosopher problem

This section illustrates the properties of augmented marked graphs obtained in the previous sections using the dining philosopher problem.

The dining philosopher problem (version 1) :

Six philosophers (H_1, H_2, H_3, H_4, H_5 and H_6) are sitting around a circular table for dinner. They are either meditating or eating the food placed at the centre of the table. There are six pieces of chopsticks (C_1, C_2, C_3, C_4, C_5 and C_6) shared by them for getting the food to eat, as shown in Figure 12. For one to get the food to eat, both the chopstick at the right hand side and the chopstick at the left hand side must be available. The philosopher then grasps both chopsticks simultaneously and then takes the food to eat. Afterwards, the chopsticks are released and returned to their original positions simultaneously.

Figure 13 shows the augmented marked graph $(N, M_0; R)$ which represents the dining philosopher problem (version 1).

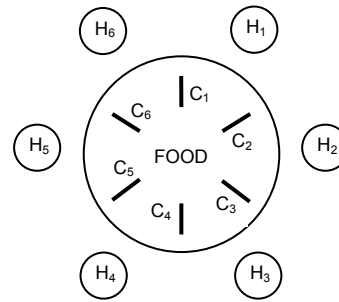


Figure 12: The dining philosopher problem.

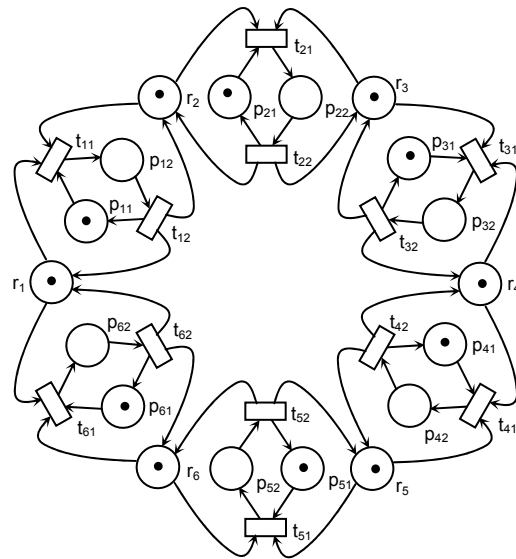


Figure 13 : The dining philosopher problem (version 1).

Semantic meaning for places and transitions

- p_{11} H_1 is meditating.
- p_{12} H_1 has got C_1 and C_2 and takes the food.
- p_{21} H_2 is meditating.
- p_{22} H_2 has got C_2 and C_3 and takes the food.
- p_{31} H_3 is meditating.
- p_{32} H_3 has got C_3 and C_4 and takes the food.
- p_{41} H_4 is meditating.
- p_{42} H_4 has got C_4 and C_5 and takes the food.
- p_{51} H_5 is meditating.
- p_{52} H_5 has got C_5 and C_6 and takes the food.
- p_{61} H_6 is meditating.
- p_{62} H_6 has got C_6 and C_1 and takes the food.
- r_1 C_1 is available for pick.
- r_2 C_2 is available for pick.
- r_3 C_3 is available for pick.
- r_4 C_4 is available for pick.
- r_5 C_5 is available for pick.
- r_6 C_6 is available for pick.
- t_{11} H_1 takes the action to grasp C_1 and C_2 .
- t_{12} H_1 takes the action to return C_1 and C_2 .
- t_{21} H_1 takes the action to grasp C_2 and C_3 .
- t_{22} H_1 takes the action to return C_2 and C_3 .
- t_{31} H_1 takes the action to grasp C_3 and C_4 .
- t_{32} H_1 takes the action to return C_3 and C_4 .
- t_{41} H_1 takes the action to grasp C_4 and C_5 .
- t_{42} H_1 takes the action to return C_4 and C_5 .
- t_{51} H_1 takes the action to grasp C_5 and C_6 .
- t_{52} H_1 takes the action to return C_5 and C_6 .
- t_{61} H_1 takes the action to grasp C_6 and C_1 .
- t_{62} H_1 takes the action to return C_6 and C_1 .

For $(N, M_0; R)$, every R-siphons contains a marked trap and would never become empty. Every place in its R-transform belongs to a cycle. Based on the results obtained in Sections 4 and 5, $(N, M_0; R)$ is live, bounded, reversible and conservative.

The dining philosopher problem (version 2) :

The Dining Philosopher Problem is revised. For one to get the food to eat, he or she first grasps the chopstick at the right hand side if available, then grasps the chopstick at the left hand side if available, and then takes the food to eat. Afterwards, the chopsticks are released and returned to their original positions simultaneously.

Figure 14 shows the augmented marked graph $(N, M_0; R)$ which represents the dining philosopher problem (version 2). The set of places $\{r_1, p_{13}, r_2, p_{23}, r_3, p_{33}, r_4, p_{43}, r_5, p_{53}, r_6, p_{63}\}$ is a R-siphon which would become empty after firing the sequence of transitions $\langle t_{11}, t_{12}, t_{13}, t_{14}, t_{15}, t_{16} \rangle$. Deadlock would occur after firing $\langle t_{11}, t_{12}, t_{13}, t_{14}, t_{15}, t_{16} \rangle$. Based on the results obtained in Section 4, $(N, M_0; R)$ is neither live nor reversible. On the other hand, for the R-transform of $(N, M_0; R)$, every place belongs to a cycle. Based on the results obtained in Section 5, $(N, M_0; R)$ is bounded and conservative.

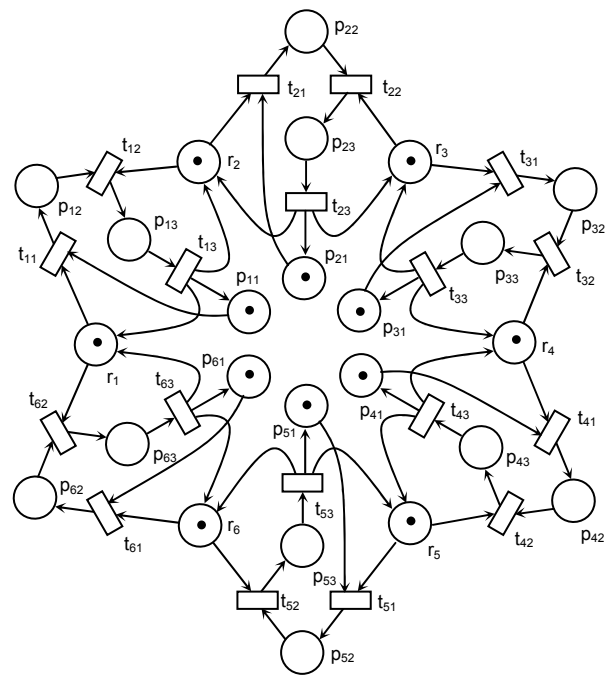


Figure 14 : The dining philosopher problem (version 2).

Semantic meaning for places and transitions

- p₁₁ H₁ is meditating.
- p₁₂ H₁ has got C₁ and prepares to pick C₂.
- p₁₃ H₁ has got C₁ and C₂ and takes the food.
- p₂₁ H₂ is meditating.
- p₂₂ H₂ has got C₂ and prepares to pick C₃.
- p₂₃ H₂ has got C₂ and C₃ and takes the food.
- p₃₁ H₃ is meditating.
- p₃₂ H₃ has got C₃ and prepares to pick C₄.
- p₃₃ H₃ has got C₃ and C₄ and takes the food.
- p₄₁ H₄ is meditating.
- p₄₂ H₄ has got C₄ and prepares to pick C₅.
- p₄₃ H₄ has got C₄ and C₅ and takes the food.
- p₅₁ H₅ is meditating.
- p₅₂ H₅ has got C₅ and prepares to pick C₆.
- p₅₃ H₅ has got C₅ and C₆ and takes the food.
- p₆₁ H₆ is meditating.
- p₆₂ H₆ has got C₆ and prepares to pick C₁.
- p₆₃ H₆ has got C₆ and C₁ and takes the food.
- r₁ C₁ is available for pick.
- r₂ C₂ is available for pick.
- r₃ C₃ is available for pick.
- r₄ C₄ is available for pick.
- r₅ C₅ is available for pick.
- r₆ C₆ is available for pick.
- t₁₁ H₁ takes the action to grasp C₁.
- t₁₂ H₁ takes the action to grasp C₂.
- t₁₃ H₁ takes the action to return C₁ and C₂.
- t₂₁ H₂ takes the action to grasp C₂.
- t₂₂ H₂ takes the action to grasp C₃.
- t₂₃ H₂ takes the action to return C₂ and C₃.
- t₃₁ H₃ takes the action to grasp C₃.
- t₃₂ H₃ takes the action to grasp C₄.
- t₃₃ H₃ takes the action to return C₃ and C₄.
- t₄₁ H₄ takes the action to grasp C₄.
- t₄₂ H₄ takes the action to grasp C₅.
- t₄₃ H₄ takes the action to return C₄ and C₅.
- t₅₁ H₅ takes the action to grasp C₅.
- t₅₂ H₅ takes the action to grasp C₆.
- t₅₃ H₅ takes the action to return C₅ and C₆.
- t₆₁ H₆ takes the action to grasp C₆.
- t₆₂ H₆ takes the action to grasp C₁.
- t₆₃ H₆ takes the action to return C₆ and C₁.

7 Conclusion

In the past decade, augmented marked graphs have evolved into a sub-class of Petri nets for modelling shared resource systems. One major reason is that augmented marked graphs possess a structure which is desirable for modelling shared resources. However, the properties of augmented marked graphs are not extensively studied.

In this paper, a number of new characterisations for live and reversible augmented marked graphs are proposed. In particular, some of these characterisations are based on cycles, instead of siphons. Besides, a R-transform is introduced. Based on the R-transform, a number of new characterisations for bounded and conservative augmented marked graphs are proposed. Consolidating these results, pretty simple conditions and procedures for checking the liveness, reversibility, boundedness and conservativeness of augmented marked graphs are derived. These have been illustrated using the dining philosopher problem.

Augmented marked graphs are often used for modelling shared-resource systems wherein the system analyst need to achieve the system design objectives on two folds. On one hand, the resources are scarce and should be maximally shared. On the other hand, the system should be carefully designed so that erroneous situations, such as deadlock and capacity overflow, due to sharing of resources should be avoided. For a shared-resource system modelled as an augmented marked graph, essential properties such as liveness, reversibility, boundedness and conservativeness can be effectively analysed with the new characterisations for augmented marked graphs. These contribute to ensuring the design correctness of shared resource systems.

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Appendix

For clarity in presentation, proofs of the proposed properties for augmented marked graphs are shown in this appendix as follows.

Proof of Property 8. (\Leftarrow) For an augmented marked graph, if every minimal siphon would never become empty, every siphon which contains at least one minimal siphon would never become empty. It follows from Properties 6 and 7 that the augmented marked graph is live and reversible. (\Rightarrow) It follows from Property 6 that every siphon (and hence, every minimal siphon) would never become empty.

Proof of Property 11. Let $S = \{ p_1, p_2, \dots, p_n \}$. For each p_i , by definition of augmented marked graphs that $\bullet p_i \neq \emptyset$. Then, there exists $p_j \in S$, where $p_j \neq p_i$, such that $(p_j \bullet \cap \bullet p_i) \neq \emptyset$. Since S is a minimal siphon, according to Property 10, p_i connects to p_j via a conflict-free path in S . Since p_j connects to p_i , this forms a cycle γ_i in S , where $p_i \in P[\gamma_i] \subseteq S$. Let $Y = \{ \gamma_1, \gamma_2, \dots, \gamma_n \}$. We have $P[Y] = P[\gamma_1] \cup P[\gamma_2] \cup \dots \cup P[\gamma_n] \subseteq S$. On the other hand, $S \subseteq (P[\gamma_1] \cup P[\gamma_2] \cup \dots \cup P[\gamma_n]) = P[Y]$ because $S = \{ p_1, p_2, \dots, p_n \}$. Hence, $P[Y] = S$.

Proof of Property 12. (by contradiction) Let $(N, M_0; R)$ be an augmented marked graph. Suppose there exists a cycle γ in $(N, M_0; R)$, such that γ is not marked. γ does not contain any place in R , and also exists in the net (N', M_0') obtained from $(N, M_0; R)$ after removing the places in R and their associated arcs. However, by definition of augmented marked graphs, γ is marked.

Proof of Property 13. For an augmented marked graph, according to Properties 11 and 12, every minimal siphon contains cycles and is marked. Hence, every siphon, which contains at least one minimal siphon, is marked.

Proof of Property 14. Let $D_r = \{ \langle t_{s1}, t_{h1} \rangle, \langle t_{s2}, t_{h2} \rangle, \dots, \langle t_{sn}, t_{hn} \rangle \}$, where $r = \{ t_{s1}, t_{s2}, \dots, t_{sn} \}$ and $\bullet r = \{ t_{h1}, t_{h2}, \dots, t_{hn} \}$. For each $\langle t_{si}, t_{hi} \rangle \in D_r$, t_{si} connects to t_{hi} via an elementary path ρ_i which is not marked. Let $S = P_1 \cup P_2 \cup \dots \cup P_n \cup \{ r \}$, where P_i is the set of places in ρ_i . We have $\bullet P_i \subseteq (P_i \bullet \cup \bullet r)$ because, for each $p \in P_i$, $| \bullet p | = | p \bullet | = 1$. Then, $(\bullet P_1 \cup \bullet P_2 \cup \dots \cup \bullet P_n) \subseteq (P_1 \bullet \cup P_2 \bullet \cup \dots \cup P_n \bullet \cup \bullet r)$. Besides, $\bullet r = \{ t_{h1}, t_{h2}, \dots, t_{hn} \} \subseteq (P_1 \bullet \cup P_2 \bullet \cup \dots \cup P_n \bullet)$. Hence, $\bullet S = (\bullet P_1 \cup \bullet P_2 \cup \dots \cup \bullet P_n \cup \bullet r) \subseteq (P_1 \bullet \cup P_2 \bullet \cup \dots \cup P_n \bullet \cup \bullet r) = S \bullet$. Therefore, S is a siphon in which r is the only one marked place. Let S' be a minimal siphon in S . According to Property 13, S' is marked. Since r is the only one marked place in S , r is also the only one marked place in S' .

Proof of Property 15. (By contradiction) Let S be a R-siphon. According to Property 11, S is covered by cycles. Suppose there exists a cycle γ in S , such that $\gamma \notin \Omega_N[R]$. By definition of augmented marked graphs, for any $p \in P[\gamma]$, $| \bullet p | = | p \bullet | = 1$. Hence, $\bullet P[\gamma] = P[\gamma] \bullet$, and $P[\gamma]$ is a siphon. Since there exists a place $r \in R$ such that $r \in S$ but $r \notin P[\gamma]$, we have $P[\gamma] \subset S$. However, since S is a minimal siphon, there does not exist any siphon $S' = P[\gamma] \subset S$.

Proof of Property 16. (by contradiction) Suppose there exists $s \in (S \setminus R)$ such that $t \in s^\bullet$. By definition of augmented marked graphs, $|s^\bullet| = |s^\circ| = 1$. S is covered by cycles in accordance with Property 15. Hence, t is the one and only one transition in s^\bullet , where $t \in T[Y] = (S^\circ \cap S)$. This however contradicts $t \in (S^\bullet \setminus S)$.

Proof of Property 17. Let S be a NR-siphon. According to Property 13, S is marked. By definition of augmented marked graphs that, for any $s \in S$, $|s^\bullet| = |s^\circ| = 1$. Then, $S = S^\bullet$ and S is also a trap. Hence, S contains itself as a marked trap and would never become empty.

Proof of Property 18. (\Leftarrow) According to Property 17, a NR-siphon would never become empty. Given that no R-siphons (and hence, no minimal siphon) eventually become empty, according to Property 8, $(N, M_0; R)$ is live and reversible. (\Rightarrow) It follows from Property 6 that no R-siphons eventually become empty.

Proof of Property 19. (\Leftarrow) According to Property 17, a NR-siphon contains a marked trap. Given that every R-siphon contains a marked trap, every minimal siphon contains a marked trap. (\Rightarrow) Since R-siphons are minimal siphons, every R-siphon contains a marked trap.

Proof of Property 20. For $(N, M_0; R)$, if every R-siphon contains a marked trap, according to Property 19, the siphon-trap property is satisfied. Hence, every minimal siphon contains a marked trap and would never become empty. It then follows from Property 8 that $(N, M_0; R)$ is live and reversible.

Proof of Property 21. Since S is a minimal siphon, according to Property 10, for any $q, q' \in S = P[Y]$, there exists in $S = P[Y]$ a conflict-free path from q to q' . Hence, Y is conflict free.

Proof of Property 22. Let $S = \{p_1, p_2, \dots, p_n\}$. According to Property 13, S is marked. It follows from Properties 15 and 21 that there exists a set of cycles $Y \subseteq \Omega_N[R]$, such that Y is conflict-free and $P[Y] = S$. Since S is a siphon, for each $p_i \in S$, $p_i^\bullet \subseteq (S^\circ \cap S^\bullet) = (P[Y] \cap P[Y]^\bullet) = T[Y]$. In case $p_i \notin R$, $p_i^\bullet \subseteq T[Y]$ because $|p_i^\bullet| = |p_i^\circ| = 1$. In case $p_i \in R$, given that p_i satisfies the R-inclusion property, $p_i^\bullet \subseteq T[Y]$. Every $p_i^\bullet \subseteq T[Y] = (P[Y] \cap P[Y]^\bullet)$ and $p_i^\bullet \subseteq P[Y] = S$. Since $S^\bullet = (p_1^\bullet \cup p_2^\bullet \cup \dots \cup p_n^\bullet) \subseteq S$, S is also a trap.

Proof of Property 23. (\Leftarrow) It follows from Properties 19 and 22. (\Rightarrow by contradiction) Suppose there exists $r \in R$, not satisfying the R-inclusion property. According to Property 14, there exists a R-siphon S , in which r is the only marked place. It follows from Properties 15 and 21 that there exists $Y \subseteq \Omega_N[R]$, such that Y is conflict-free and $S = P[Y]$. According to Property 19, S contains a marked trap Q . Then, $r \in Q$ and $r^\bullet \subseteq (Q^\circ \cap Q^\bullet)$. Since S is a siphon, we have $r^\bullet \subseteq (S^\circ \cap S^\bullet) = (P[Y] \cap P[Y]^\bullet) = T[Y]$. However, as r does not satisfy the R-inclusion property, $r^\bullet \not\subseteq T[Y] = (P[Y] \cap P[Y]^\bullet) = (S^\circ \cap S^\bullet)$, implying $r^\bullet \not\subseteq (Q^\circ \cap Q^\bullet)$.

Proof of Property 24. According to Property 23, $(N, M_0; R)$ satisfies the siphon-trap property. It follows from Property 20 that $(N, M_0; R)$ is live and reversible.

Proof of Property 25. For each place $p \notin R$ in N , $M_0; R$, $|p^\bullet| = |p^\circ| = 1$. Each place $r \in R$ is replaced by a set of places $\{p_1, p_2, \dots, p_{k_r}\}$, where $|p_i^\bullet| = |p_i^\circ| = 1$ for $i = 1, 2, \dots, k_r$. Hence, for every place q in N' , $|q^\bullet| = |q^\circ| = 1$. (N', M_0') is a marked graph.

Proof of Property 26. Let (N', M_0') be the R-transform of an augmented marked graph $(N, M_0; R)$. As the transformation does not create cycles, cycles in (N', M_0') exist in $(N, M_0; R)$. According to Property 12, cycles in $(N, M_0; R)$ are marked, and hence, cycles in (N', M_0') are marked. Since (N', M_0') is a marked graph, it follows from Property 2 that (N', M_0') is live.

Proof of Property 27. Let $D_r = \{\langle t_{s_1}, t_{h_1} \rangle, \langle t_{s_2}, t_{h_2} \rangle, \dots, \langle t_{s_{k_r}}, t_{h_{k_r}} \rangle\}$. By definition of augmented marked graphs, for each $\langle t_{s_i}, t_{h_i} \rangle$, there exists an unmarked path $\rho = \langle t_{s_1}, \dots, t_{h_1} \rangle$ in $(N, M_0; R)$. Hence, ρ also exists as an unmarked path in (N', M_0') , and ρ together with q_i forms a cycle γ_i which is marked at q_i only. Since (N', M_0') is a marked graph, according to Property 5, the corresponding vector of γ_i is a place invariant α_i of N' . Since q_i is the only one marked place in γ_i , $\alpha_i[q_i] = 1$ and $\alpha_i[s] = 0$ for any $s \in P_0 \setminus \{q_i\}$.

Proof of Property 28. Let $P = \{p_1, p_2, \dots, p_n\}$ be the places in N' , and $P_0 \subseteq P$ be those marked places. Since each p_i belongs to a cycle γ_i and (N', M_0') is a marked graph, according to Property 5, the corresponding vector of γ_i is a place invariant α_i' of N' . Then, $\alpha' = \alpha_1' + \alpha_2' + \dots + \alpha_n' > 0$ is a place invariant of N' . Consider $Q_i = \{q_1, q_2, \dots, q_k\}$. Let $q_m \in Q_i$ such that $\alpha'[q_m] \geq \alpha'[q_j]$ for any $q_j \in Q_i$. For each q_j , according to Property 27, there exists a place invariant $\alpha_j' > 0$ such that $\alpha_j'[q_j] = 1$ and $\alpha_j'[s] = 0$ for any $s \in P_0 \setminus \{q_j\}$. There also exists a place invariant $\alpha'' = \alpha' + h\alpha_j'$, where $h \geq 1$, such that $\alpha''[q_j] = \alpha'[q_m]$ and $\alpha''[s] = \alpha'[s]$ for any $s \in P_0 \setminus \{q_j\}$. Therefore, there eventually exists a place invariant α of N' such that $\alpha[q_1] = \alpha[q_2] = \dots = \alpha[q_k]$.

Proof of Property 29. (\Leftarrow) Let $R = \{r_1, r_2, \dots, r_n\}$, where each r_i is being replaced by a set of places Q_i , for $i = 1, 2, \dots, n$. Since every place in (N', M_0') belongs to a cycle, according to Property 28, there exists a place invariant α' of N' such that $\alpha' > 0$ and $\alpha'[q_1] = \alpha'[q_2] = \dots = \alpha'[q_k]$ for each $Q_i = \{q_1, q_2, \dots, q_k\}$. Intuitively, there also exists a place invariants α of N such that $\alpha > 0$ and $\alpha[r_i] = \alpha'[q_1] = \alpha'[q_2] = \dots = \alpha'[q_k]$ for each Q_i . Hence, $(N, M_0; R)$ is conservative. According to Property 1, $(N, M_0; R)$ is also bounded. (\Rightarrow) Since $(N, M_0; R)$ is conservative, there exists a place invariant α of N such that $\alpha > 0$. Consider each $r_i \in R$ which is being replaced by $Q_i = \{q_1, q_2, \dots, q_k\}$. Intuitively, there also exists a place invariant α' of N' such that $\alpha' > 0$ and $\alpha'[q_1] = \alpha'[q_2] = \dots = \alpha'[q_k] = \alpha[r_i]$ and $\alpha'[s] = \alpha[s]$ for any $s \in P \setminus Q_i$. Hence, (N', M_0') is conservative. It follows from Property 1 that (N', M_0') is also bounded. Since (N', M_0') is a marked graph, according to Property 3, every place in (N', M_0') belongs to a cycle.

Proof of Property 30. It follows from Properties 3 and 29.